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1D Grid of 1D Blocks Cuda Kernels A CUDA Kernel function is defined using the __global_ keyword. 1D Grid of 1D Blocks A Kernel is executed N times in parallel by N different threads on the Grid device Each thread has a unique ID stored in the built-in threadIdx variable, Block 0 Block 1 Block 2 a struct with components x,y,z. Threads Threads Threads Each thread block has a unique ID stored in the built-in *blockIdx* 0 1 2 3 0 1 2 3 0 1 2 3 variable, a struct with components x,y,z. int index = blockIdx.x * blockDim.x + thread Idx.x Kernel Configuration kernel Fun cti on< <<n um blocks, num th reads> Kernel Execution >(params) 1D Grid of 3D Blocks Configuration 1D Grid of 3D Blocks num_blocks The number of thread blocks along each dimension of the grid. Grid The number of threads along each dimension of the thread block num th-Block 0 Block 1 reads 0,1 1,0,1 2,0 1,0,0 2,0,0 1,0,0 2,0,0 ooŏ 000 **CUDA Thread Organization** Threads Threads Thread are grouped in blocks and can be organized in 1 to 3 dimens-1,1,0 2, 1,1,0 2 ions. Blocks are grouped into grids which can be organized in 1 to 3 dimensions. int index = blockIdx.x blockDim.x blockDim.y bl Blocks are executed independently. ockDim.z + thread Idx.z blockDim.y blockDim.x + thread Idx.y blockDim.x + thread Idx.x; 2D Grid of 2D Blocks applied on a Matrix MatrixWidth (2,1 The index of each thread is identified by two coordinates i and j. We can find i applying the rule of 1D Grid of 1D Blocks over the x axis: int i = blockIdx.x * blockDim.x + thread Idx.x; And we can find j applying the rule of 1D Grid of 1D Blocks over the y axis: int j = blockIdx.y * blockDim.y + thread Idx.y; Thus, knowing that a row in the grid is large *GridDim.x times* BlockDim.x, we can calculate the index: int index = j gridDim.x blockDim.x +i; By m_amendola Published 22nd July, 2023.

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CUDA Events

Declaring a Cuda Event	cudaEv ent_t event;
Allocating the event	cudaEv ent Cre ate (& event);
Recording the Event.	cudaEv ent Rec ord (ev ent);
Synchronizing the event	cudaEv ent Syn chr oni ze(event);
Find elapsed time between two events	cudaEv ent Ela pse dTi me(&e lapsed, a, b);
Free event variables	cudaEv ent Des tro y(e vent);

CUDA Streams

GPU operations on CUDA use execution queues called streams.

Operations pushed in a stream are executed according to a FIFO policy.

There is a default Stream, called stream 0.

Operations pushed in a non-default stream will be executed after all operations on default stream are emptied.

Operations assigned to default stream introduce implicit synchronization barriers among other streams.

CUDA Streams API

Create a streamcudaSt rea mCr eat e(s tream1);Deallocate a streamcudaSt rea mDe str oy(stream)



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CUDA Streams API (cont)

Block host until all	cudaSt	rea	mSy	nch	ron	ize	(st
operations on a stream	ream)	;					
are completed.							

We can use stream to obtain the concurrent execution of the same kernel or different kernels.

Synchronization operations

Synchronization operations	
Explicit Synchronization	Implicit Synchr- onization
<i>cudaDeviceSynchronize()</i> blocks host code until all operations on device are completed	Operations assigned to default stream
<i>cudaStreamWaitEvent(stream, event)</i> blocks all operations assigned to a stream until event is reached.	Memory Alloca- tions on device
	Settings operations on device
	Page-locked memory alloca- tions

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CUDA API			Memory Allocation AP	I Functions (cont)
https://docs.nvidia.com/cuda/cuda-runtime-api/index.html			Cp y(void <i>dst, void</i> src, size_t size	
Then we at to the alloc Finally, we memory fro At the end the device Copy opera	orkflow ocate and "build" the input on the hos llocate dynamic memory on the device ated memory areas. initialize the memory on the device a om the host to the device. of the computation, we may want to o to the host. ation is <i>blocking</i> .	e , obtaining pointers and we copy the	data vice); from host to device	
Dynamic memory allocation	cudaMalloc ((void **) &ude	v, N*size of(dou b	the pointer to the	Cp y(void <i>dst, void</i> src, size_t size
Memory Initializ- ation on device	<pre>cudaMe mse t(void *devPtr, t;</pre>	int val, size_t co	Jou from devPtr is device pointer to hotst the After 4.0, CUDA support After 4.0, CUDA support the systems itself known parameter The Global Memory mills the Declaring a static mills the Declaring a static variable Variable Deall@Cattofy a dynafficavariable with the constant byte value	<pre>orts Unified Virtual Addressing meaning that ws where the buffer is allocated. The direction t to cudaMemcpyDefault. </pre>
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Global Memory (cont)	Texture Memory
Allocating cudaMa llo cPi tch (&ptr, &p itch, width* siz	Managing texture memory
an , height)	Allocate cudaMa llo c(&M, memsize)
aligned 2D buffer	global
where	memory
elements	on device
are	Create a textur e <d ata="" dim="" type,=""> Mtextu reRef;</d>
padded	texture reference.
so that	
each row	Create a cudaCh ann elF orm atDesc Mdesc = cudaCr e channel tat vpe >():
is aligned	channel tat ype >(); descriptor
cudaMallocPitch returns an integer pitch that can be used to access	Bind the cudaBi ndT ext ure(0, Mtextu reRef, M, Mde
row element with stride access. For example:	texture
<pre>float *row = devPtr + r * pitch;</pre>	reference
Charad Mamon	to
Shared Memory	memory.
Static variable declar	Unbind at cudaUn bin dTe xtu re(MTe xtu reRef);
ation inside the kernel.	the end.
Dynamic variable externshared type *shmem;	<pre>In order to text1D fet ch(Mte xtu reRef, address);</pre>
allocation outside the kernel	access
	the
Constant memory	texture
Declaring cons tant type variab lename;	memory, we can
a static	use the
variable	texture
Copy cudaMe mcp yTo Sym bol (va ria ble name, &h	oseference, sizeof (type), cudaMe mcp vH
memory o stT oDe vice);	Mtextu-
from host	reRef.*
to	Accessing text2D fet ch(Mte xtu reRef, address);
device.	2D cuda
We cannot declare a dynamic variable on the costant memory	array.
	Accessing text3D fet ch(Mte xtu reRef, address);
	3D cuda
	array.



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Asynchrono	ous Data Transfers	Error Handling			
Allocates page-l- ocked memory on the host.	cudaMa llo cHo st(buffer, size)	All CUDA API functions returns an error code of type <i>cudaError</i> . The constant <i>cudaSuccess</i> means no error. <i>cudaGetLastError</i> return the status of the internal error variable. Calling this function resets the internal error to cudaSuccess.			
Frees page-l- ocked memory.	cudaFr eeH ost (bu ffer)	<pre>Macro for Error Handling #define CUDA_CHECK(X) {\ cudaEr ror_t _m_cud aStat = X; \ if(cud aSu ccess != _m_cud aStat) {\</pre>			
Registers an existing host memory range for use by CUDA.	cudaHo stR egi ster()	<pre>fprint f(s tde rr, " \nC UDA _ERROR: %s in file %s line %d\n",\ cudaGe tEr ror Str ing (_m _cu daS tat),FILE,LINE);\ exit(1);\ } } CUDA_C HECK(cudaMe mcp y(d _buf, h_buf, buffSize,</pre>			
Unregi- sters a memory range that was registered with cudaHo- stRegi- ster.	cudaHo stU nre gis ter()	cudaMe mcp yHo stT oDe vice));			
Copies data between host and device.	cudaMe mcp yAs ync (de st_ buffer, src_bu ffe ,st ream)	er, dest_size, src_size, direct ion			
These oper	rations must be queued into a non-default stream.				

Page-locked Memory

Pageable memory is memory which is allowed to be paged in or paged out whereas **page-locked memory** is memory not allowed to be paged in or paged out.

Page out is moving data from RAM to HDD, while *page in* means moving data from HDD to RAM. These operations occurs when the main memory does not have enough free space.

Source: https://leimao.github.io/blog/Page-Locked-Host-Memory--Data-Transfer/



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