

Pointers

Storing, data type of pointers and variables must be the same

`&var` Returns address of memory location of variable

`data_type` Initialize. Put * in front of name

`*pointer;`

`*pointer` Returns value at the memory location stored by pointer

Array variables are actually pointers to the first element in the array. The amount that your pointer moves from an arithmetic operation

`*array` Returns first element in array

`*(array+2)` Returns third element in array

⊙ Variables that you declare are stored in a memory location in your computer

⊙ The address of these memory locations can be stored in pointers

⊙ Addresses are in hexadecimal

Iterators

//Append `::iterator` to your data type declaration to create an iterator of

that data type

`vector<int>::iterator it;` // declares an iterator of `vector<int>`

// loops from the start to end of `vi`

`for(vector<int>::iterator it = vi.begin(); it != vi.end(); it++)`

`cout << *it << " ";` // outputs 1 2 3 4

`deque<int> d;`

`deque<int>::iterator it;`

`it = d.begin();` //Points to first element

`it++;` // Points to next element

`it = d.end();` // Points to Last element

`it--;` // Points to previous element

`cout << *it;` // outputs the element it is pointing

Iterators are essentially pointers to an STL data structure

Maps

`map<string, int> M;`

`M["hello"] = 2;`

`M["asd"] = 986;`

`M.count("asd");` // returns 1

`M.count("doesnt_exist");` // returns 0

`M.size();`

// Check for the existence of some key in the map - $O(\log N)$

`it = M.find("asd");` //returns the iterator to "asd"

`it = M.upper_bound("aaa");`

`it = M.lower_bound("aaa");`

`if (it == M.end())`

`cout << "Does not exist\n";`

//Iteration

`for (auto it = mp.begin(); it != mp.end(); ++it) {`
`cout << it.first << ", " << it.second << "\n";`
`}`

A data structure that takes in any data[key]

Gives you the associated value stored through $O(\log N)$ magic

Best used when you need to lookup certain values in $O(\log N)$ time that are associated with other values

Queue

`queue<int> q;`

`q.push(5);` // Inserts/ Enqueue element at the back of the queue

`q.front();` // Returns element atb the front of the queue

`q.pop` // Removes (Dequeues) element from the front of the queue

`q.empty();` // Returns boolean value of whether queue is empty

First In First Out data structure where elements can only be added to the back and accessed at the front

Priority Queue

```
priority_queue<data_type> pq; // Largest at top
priority_queue<data_type, vector<data_type>,
greater<data_type> >
pq; // Smallest at top
pq.push(5); // pushes element into it. Duplicates
are allowed
pq.top() // Returns largest or smallest element
pq.pop() // Removes largest or smallest element
pq.size(); // Returns size
pq.empty(); Check if queue is empty
```

Like a queue except that only the element with the greatest priority (eg. largest/smallest) can be accessed

Fenwick tree

```
//Below here can mix & match
long long ft[100001]; // note: this fenwick tree
is 1-indexed.
////PU-
PQ////////////////////////////////////////
/////
void fenwick_update(int pos, long long value) {
    while (pos <= N) {
        //cout<<"Fenwick Updating: "<<pos<<" "<<-
<value<<endl;
        ft[pos] += value;

        pos += pos&-pos;
    }
}

long long fenwick_query(int pos) {
    long long sum = 0;
    while (pos) { // while p > 0
        sum += ft[pos];
        pos -= pos&-pos;
    }
    return sum;
}

////RU-
PQ////////////////////////////////////////
/////
void fenwick_range_update(int pos_a, int pos_b, int
val){
    //TLE way
```

Fenwick tree (cont)

```
//for (int i=pos_a;i<=pos_b;i++){fenwick_upd-
ate(i, val);}
    fenwick_update(pos_a, val);
    fenwick_update(pos_b+1, -val);
}
////PURQ////////////////////////////////
////////////////////////////////
long long fenwick_range_query(int pos_a, int
pos_b) {
    return fenwick_query(pos_b) - fenwick_query(p-
os_a-1);
}
////RU-
RQ////////////////////////////////////////
/////
long long B1[100001];long long B2[100001];
void base_update(long long *ft, int pos, long long
value){
    //Add largest power of 2 dividing x / Last set
bit in number x
    for (; pos <= N; pos += pos&(-pos))
        ft[pos] += value;
}
void rurq_range_update(int a, int b,long long v){
    base_update(B1, a, v);
    base_update(B1, b + 1, -v);
    base_update(B2, a, v * (a-1));
    base_update(B2, b + 1, -v * b);
}
void rurq_point_update(int a, long long v){
    rurq_range_update(a,a,v);
}
long long base_query(long long *ft,int b){
    long long sum = 0;
    for(; b > 0; b -= b&(-b))
        sum += ft[b];
    return sum;
}
// Return sum A[1...b]
long long rurq_query(int b){
```



Fenwick tree (cont)

```
return base_query(B1, b) * b - base_query(B2, b);
}
//Return sum A[a..b]
long long rurq_range_query(int a,int b){
return rurq_query(b) - rurq_query(a-1);
}
```

Pair

```
// Initialise
pair<data_type_1, data_type_2> variable;
// OR
pair<data_type_1, data_type_2> variable = make_p-
air(value1,value2);
// Store values
variable.first = value;
variable.second = value;
// Retrieve values
cout <<variable.first << " " << variable.second <<
endl;
//Nesting pairs
pair<int, pair<int, int> > a;
a.first = 5;
a.second.first = 6;
a.second.second = 7;
```

Stores a pair of values

Stack

```
stack<int> s;
s.push(5); // push an element onto the stack -
O(1)
s.pop(); // pop an element from the stack - O(1)
s.top(); // access the element at the top of the
stack - O(1)
s.empty(); // whether stack is empty - O(1)
```

First-In-Last-Out data structure

Only Element at the top can be accessed / removed

Vector

```
// Initialize
vector<data_type> v;
v.push_back(value); // Add element to back
v.pop_back() // Remove last element
v.clear(); // Remove all elements
v[index] // Return element of index
v.back(); // Return last element
v.size(); // Return Size of vector
v.empty() // Return boolean value of whether
vector is empty
```

Like arrays but re-sizable. You can add and remove any number of elements from any position.

Sets and Multisets

```
set<int> s; set<int>::iterator it;
multiset<int> s; multiset<int>::iterator it;
s.insert(10);
it = s.find(8)
it = s.upper_bound(7);
it = s.lower_bound(7);
s.erase(10); //Remove element from set
s.erase(it) //Can also use iterators
s.empty();
s.clear();
// to loop through a set
for(it = s.begin(); it != s.end(); it++)
cout << *it << " "; // outputs 2 7 10
```

In a set: All elements are sorted and no duplicates

Multisets can store duplicates though



Deque

```
deque<int> d;
// access an element / modify an element (0-indexed
as well) - O(1)
d[0] = 2; // change deque from {5, 10} to {2, 10}
d[0]; // returns a value of 2
d.back(); // get back (last) element - O(1)
d.front(); // get front (first) element - O(1)
d.clear() // Remove all elements
d.push_back(5); // add an element to the back -
O(1)
d.push_front(10); // add an element to the front -
O(1)
d.pop_back(); // remove the back (last) element -
O(1)
d.pop_front(); // remove the front (first) element
- O(1)
d.size(); //Return size
d.empty // Whether queue is empty
```

A stack and queue combined.

...or a vector that and be pushed and popped from the front as well.

Deque = Double Ended Queue!

Segment Tree

```
struct node {
    int start, end, mid, val, lazyadd;
    node left, right;

    node(int _s, int _e) {
        //Range of values stored
        start = _s; end = _e; mid = (start+end)/2;
        //Min value stored
        val = 0; lazyadd = 0;
        if (start!=end) {
            left = new node(start,mid);
            right = new node(mid+1,end);
        }
    }

    int value(){
        if (start==end){
            val += lazyadd;lazyadd = 0;return val;
```

Segment Tree (cont)

```
        }else{
            val += lazyadd;
            // Propagate Lazyadd
            right->lazyadd += lazyadd;
            left->lazyadd += lazyadd;
            lazyadd = 0;
            return val;
        }
    }

    void addRange(int lower_bound, int upper_
bound, int val_to_add){
        if (start == lower_bound && end == upper_
bound){
            lazyadd += val_to_add;
        }else{
            if (lower_bound > mid){
                right->addRange(lower_bound,
upper_bound, val_to_add);
            }else if (upper_bound <= mid){
                left->addRange(lower_bound,
upper_bound, val_to_add);
            }else{
                left->addRange(lower_bound, mid,
val_to_add);
                right->addRange(mid+1, upper_
bound, val_to_add);
            }
            val = min(left->value(), right->va-
lue());
        }
    }

    // Update position to new_value // O(log N)
    void update(int pos, int new_val) { //position
x to new value
        if (start==end) { val=new_val; return; }
        if (pos>mid) right->update(pos, new_val);
        if (pos<=mid) left->update(pos, new_val);
        val = min(left->val, right->val);
    }

    // Range Minimum Query // O(log N)
```

Segment Tree (cont)

```

int rangeMinimumQuery(int lower_bound, int
upper_bound) {
    //cout<<"Node:"<<start<<" "<<end<<" "<<m-
id<<" "<<val<<endl;
    //If Query Range Correspon-
ds////////////////////////////////
    if (start==lower_bound && end==upper_b-
ound) {
        return value();
    }
    //Query Right Tree if range only lies
there
    else if (lower_bound > mid) {
        return right->rangeMinimumQuery(lower-
_bound, upper_bound);
    }
    //Query Left Tree if range only lies there
    else if (upper_bound <= mid) {
        return left->rangeMinimumQuery(lower_-
bound, upper_bound);
    }
    //Query both ranges as range spans both
trees
    else {
        return min(left->rangeMinimumQuery(lo-
wer_bound, mid),right->rangeMinimumQuery(mid+1,
upper_bound));
    }
    //E-
nd////////////////////////////////////////
}

} *root;
void init(int N) {
    root = new node(0, N-1); // creates seg tree
of size n
}
void update(int P, int V) {
    root->update(P,V);
}
int query(int A, int B) {
    int val = root->rangeMinimumQuery(A,B);
    return val;
}

```

