

### Algoritmi di ordinamento

ALGORITMO	Tempo	Spazio
Selection Sort	$\Theta(n^2)$	sul posto $\Theta(1)$
Insertion Sort *	Ottimo: $\Theta(n)$ Medio: $\Theta(n^2)$ Pessimo $\Theta(n^2)$	sul posto $\Theta(1)$
MergeSort *	$\Theta(n \log n)$	$\Theta(n)$
QuickSort	Ottimo $\Theta(n \log n)$ Medio: $\Theta(n \log n)$ Pessimo: $\Theta(n^2)$	sul posto $\Theta(1)$ (ma spazio non costante x ricorsione)
HeapSort	$\Theta(n \log n)$	sul posto $\Theta(1)$
CountingSort *	$\Theta(n + k)$	
RadixSort *	$\Theta(d(n+k))$	

### SelectionSort

```

for i = 1 to n-1 {
  minimo = i;
  for j=i+1 to n {
    if (a[j] < a[minimo]) minimo = j;
  }
  scambia a[j] e a[minimo];
}

```

### Insertion Sort

```

for j = 2 to n {
  key = a[j];
  i = j - 1;
  while (i > 0 && a[i] > key) {
    a[i+1] = a[i];
    i--;
  }
  a[i+1] = key;
}

```

### MergeSort

```

MergeSort(a, sx, dx)
if (sx < dx) {
  cx = (sx+dx)/2; //parte intera inferiore
  MergeSort(a, sx, cx);
  MergeSort(a, cx+1, dx);
  Merge(a, sx, cx, dx);
}
Merge(a, sx, cx, dx)
n1 = cx - sx + 1;
n2 = dx - cx;
L = nuovo array di dim n1+1
R = nuovo array di dim n2+1
for(i = 1; i ≤ n1; i++) L[i] = a[sx+i-1];
for(j = 1; j ≤ n2; j++) R[j] = a[cx+j];
L[n1+1] = ∞ //sentinella
R[n2+1] = ∞ //sentinella
i = 1;
j = 1;
for(k=sx; k ≤ dx; k++) {
  if (L[i] ≤ R[j]) {a[k] = L[i]; i++;}
  else {a[k] = R[j]; j++;}
}

```

### QuickSort

```

QuickSort(A, p, r)
if (p < r) {
  q = partition(A, p, r);
  QuickSort(A, p, q-1);
  QuickSort(A, q+1, r);
}
Partition(A, p, r) {
  //nella versione random generare i = numero
  //random tra p e r
  //e metterlo in fondo
  x = A[r];
  i = p-1;
  for j = p to r-1 {
    if(a[j] ≤ x) {
      i++;
      scambia A[i] con A[j];
    }
  }
  scambia(A[i+1], A[r]);
  return i+1;
}

```

### Heap

```

Max-Heapify(A, i)
l = Left(i);
r = Right(i);
if l ≤ A.heapsize and A[l] > A[i]
  massimo = l;
else massimo = i;
if r ≤ A.heapsize and A[r] > A[massimo]
  massimo = r;
if massimo ≠ i
  scambia A[i] con A[massimo];
  Max-Heapify(A, massimo);
Build-Max-Heap
A.heapsize = A.length;
for i = A.length/2 downto 1
  Max-Heapify(A, i);
Heapsort(A)
Build-Max-Heap(A);
for i = A.length downto 2
  scambia A[1] con A[i]
  A.heapsize--;
  Max-Heapify(A, 1);

```

### Code di priorità

```

Heap-Maximum(A) return A[1];
Heap-Extract-Max(A)
if A.heap-size < 1
  error "underflow dell'heap"
max = A[1];
A[1] = A[A.heapsize]
A.heapsize--;
Max-Heapify(A, 1);
return max;
Heap-Increase-Key (A, i, key)
if key < A[i]
  error "la nuova chiave è più piccola di quella corrente";
A[i] = key;
while i > 1 and A[Parent(i)] < A[i]
  scambia A[i] con A[Parent(i)];
  i = Parent(i);
Max-Heap-Insert(A, key)
A.heapsize++;
A[A.heapsize] = -∞;
Heap-Increase-key (A, A.heapsize, key)

```



By **greatcraic**  
[cheatography.com/greatcraic/](https://cheatography.com/greatcraic/)

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### CountingSort

```
CountingSort(A, k)
C[0..k] nuovo array di dimensione k
//nota che qui conto da 0
for i = 0 to k
  C[i] = 0;
for i = 0 to n
  C[A[i]]++;
for i = 1 to k
  C[i] = C[i-1] + C[i];
for i = n downto 1 {
  B[C[A[i]]] = A[i];
  C[A[i]]--;
}
```

### RadixSort

```
RadixSort(A, d)
for i = 1 to d
  usa un ordinamento stabile per ordinare
  l'array A sulla cifra i
```

### Ricerca Binaria (tempo log n)

```
RicercaBinaria(a, sx, dx, k)
if (sx > dx) return -1;
cx = sx + dx / 2;
if (a[cx] == k) return cx;
if (a[cx] > k)
  return RicercaBinaria(a, sx, cx-1, k);
else
  return RicercaBinaria(a, cx+1, dx, k);
RicercaBinariaSx(a, sx, dx, k)
if (sx > dx) return -1;
if (sx == dx) {
  if (a[sx] == k) return sx;
  else return -1;
}
cx = sx + dx / 2;
if (a[cx] ≥ k)
  return RicercaBinaria(a, sx, cx, k);
else
  return RicercaBinaria(a, cx+1, dx, k);
```

### Randomized-Select (QuickSelect)

```
//Trova la i-esima occorrenza + piccola in
//tempo atteso lineare
Randomized-Select(A, p, r, i)
if (p == r)
  return A[p];
q = Randomized-Partition(A, p, r); //vedi
//quicksort
k = q - p + 1;
if (i == k) // il valore del pivot è la
  // soluzione
  return A[q];
else if (i < k)
  return Randomized-Select(a, p, q-1, i)
else
  return Randomized-Select(a, q+1, r, i-k)
```



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