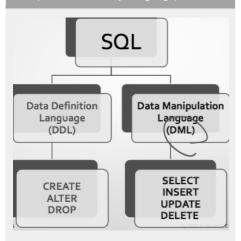


CS 411 Midterm 1 Cheat Sheet

by etau97hi1 via cheatography.com/32178/cs/18949/

SQL (Structured Query Language)



Relational Database Model

- Data is structured as relations
- Defines a limited set of **operations** (query and modification) to interact with the data
- Allow for defining **constraints** on columns (attributes), a table (relation), relationships among tables (foreign keys).

Syntax: Patterns and "LIKE"

- Used in a WHERE clause
- Pattern: quoted string with % for any string, _
 for any character

Closed

Query language is **closed** if we can use the answer from one query as input to another query

SQL Subquery Example

SELECT *

FROM (SELECT * FROM Customer WHERE
 name LIKE 'A%') as temp
WHERE temp.phone LIKE '5%';

By etau97hi1

cheatography.com/etau97hi1/

SQL Subqueries that Return Scalar

Subquery can be used as value if guaranteed to produce one tuple with one component

- "Single" tuple often guaranteed by key constraint
- A run-time error if not scalar

SQL: Boolean Operators

<tuple> IN <relation> is true if and only if the
tuple is a member of the relation.

EXISTS(<relation>) is true if and only if the <relation> is not empty.

 $x = ANY(\mbox{\ensuremath{<}}\mbox{\ensuremath{<}}\mbox{\ensuremath{<}}\mbox{\ensuremath{>}}\mbox{\ensuremath{<}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}}\mbox{\ensuremath{>}$

Similarly, $x \leftrightarrow ALL(< relation >)$ is true iff for every tuple t in the relation, x is not equal to t

Ex: Aggregations

SELECT AVG(price) FROM Sells
WHERE drink = 'Mocha';

Aggregations

SUM, AVG, COUNT, MIN, and MAX can be applied to a column in a SELECT clause to produce that aggregation on the column.

Ex: Grouping

SELECT customer, AVG(price)

FROM Frequents, Sells

WHERE drink = 'Mocha' AND

Frequents.cafe = Sells.cafe

GROUP BY customer

Single User Assumptions:

- Each operation (UPDATE ... SET ... WHERE) is executed one at a time
- **ISOLATION** one op exec, maybe change DB, then next op exec
- ATOMIC op exec entirely or not at all

Transactions Definition

- Group the SFW and USW into atransaction
- Transaction is a sequence of statements considered a "unit of operation" on DB
- Serializability of transactions Either user1 transaction exec first or user2's, but not in parallel

Transactions

Transaction: sequence of read/write ops on the DB w/ the property that either all or none of actions complete.

- May either succeed (COMMIT), or fail (ABORT or ROLLBACK)

ACID Properties

Atomicity either all ops exec or none

Consiste trans exec in isolation keeps DB in consistent state

Isolation trans isolated from effects of other concurrently exec trans

Durability updates stay in DBMS

Transaction Manager

Ensure that transactions that exec in parallel don't interfere with each other.

Concurrent Execution Problems

Write-Read conflict dirty/inconsistent read
Read-Write conflict Unrepeatable read
Write-Write conflict lost update

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