

Search Methods

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|--------------|---|
| Tree Search | Expand nodes using gradients |
| Graph Search | Avoids revisiting nodes and ensure efficiency |

Uninformed search

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| Uniform cost search | aka Cheapest-first Add visited node to Explored and add its neighbors to the frontier Visit cheapest node in the frontier Move to next cheapest if all neighbors are explored |
| Iterative deepening | Iteratively calls depth-limited search Initialize frontier with root If not goal, remove from frontier and expand If at the end of depth, terminate and start over Repeat until goal is found Guaranteed to find optimal path More efficient than DFS Time complexity: $O(b^d)$ Space complexity: $O(d)$ |
| Bidirectional | Finds shortest path of a graph |

Informed Search

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|--------------------------|--|
| Best-first search | Choose unvisited node with the best heuristic value to visit next Same as lowest cost BFS |
| Greedy best-first search | Uses heuristic to get the node closest to the goal Bad performance if heuristic is bad Does NOT consider cost of getting to the node |

Informed Search (cont)

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|----------------------|--|
| A* search | Always expand to node with minimum f Evaluate cost of getting to goal using heuristics $f(n) = g(n) + h(n)$ where g is cost to get to n Uses priority queue |
| Heuristics | Cost to get to the goal |
| Admissible heuristic | Optimistic model for estimating cost to reach the goal Never overestimates $h(n) \leq c(n)$ where c is actual cost |
| Consistent heuristic | $h(n) \leq c(n, a, n') + h(n')$ Immediate path costs less than longer path Consistent \implies Admissible |

Consistent heuristic

Consistent heuristics

- A heuristic is consistent if for every node n , every successor n' of n generated by any action a ,
$$h(n) \leq c(n, a, n') + h(n')$$
- If h is consistent, we have
$$\begin{aligned} f(n) &= g(n) + h(n) \\ &= g(n) + c(n, a, n') + h(n') \\ &\geq g(n') + h(n') \\ &= f(n') \end{aligned}$$
- i.e., $f(n)$ is non-decreasing along any path.
- Theorem: If $h(n)$ is consistent, A* using GRAPH-SEARCH is optimal



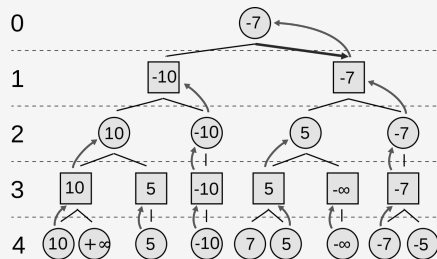
Adversarial Search

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| Hill climbing | Method of local search Only move to neighbors to find the max Does NOT guarantee to find optimal |
| Simulated annealing | Method of local search Combine hill climbing and random walk Always find global max |

Adversarial Search (cont)

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|-------------------|---|
| Local beam | Generate k random states Generate successors of all k states If goal stop; else, pick k best successors and repeat Different from hill-climbing since information is shared between k points |
| Genetic algorithm | Cross-Over and mutation Decomposes strands of DNA and permute Produces children by: Selection, Cross-over, Mutation |

Minimax Tree



⊙ Max node
□ Min node

α-β Pruning

```
function alphabeta(node, depth, α, β,
maximizingPlayer)
    if depth = 0 or node is a terminal node
        return the heuristic value of node
    if maximizingPlayer
        v := -∞
        for each child of node
            v := max(v, alphabeta (child,
depth - 1, α, β, FALSE))
            α := max(α, v)
            if β ≤ α
                β cutoff
            return v
    else
        v := ∞
        for each child of node
            v := min(v, alphabeta (child,
depth - 1, α, β, TRUE))
            β := min(β, v)
            if β ≤ α
                α cutoff
            return v
```

SAT

| P | Q | ¬P | P∨Q | P∧Q | P⇒Q | ¬P∨Q |
|-------|-------|-------|-------|-------|-------|-------|
| False | False | True | False | False | True | True |
| False | True | True | True | False | True | True |
| True | False | False | False | False | False | False |
| True | True | False | True | True | True | True |

Propositional SAT: Graph coloring

| | |
|--|---|
| At least 1 of K per i | $(C_i, 1 \vee C_{i,2} \vee \dots \vee C_{i,k})$ $O(n)$ clauses |
| $1 \geq \text{color per } i$ | $\forall k \neq k' (\neg C_{i,k} \vee \neg C_{i,k'})$ $O(n^2)$ |
| If node i and j share an edge assign different colors | $\forall x \in E, (\neg C_{i,x} \vee \neg C_{j,x})$ |